

Conserving Power in MANETs using Framework

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Abstract

Background/Objective: This paper presents a new digital control approach to track the Maximum Power Point (MPPT) for Photovoltaic (PV) systems. A photovoltaic power generation system requires an effective controller to overcome sudden changes to irradiance and to maximize its efficiency. **Methods/Statistical Analysis:** A new approach to MPPT based on voltage control was proposed, while a PI (Proportional Integral) controller was combined with the boost converter to adapt the duty cycle. **Findings:** The input voltage reference was adaptively perturbed with variable steps until the maximum power was reached. A state-space model was derived through the averaging method, with the control input being the duty ratio to regulate the pulse width modulator on the DC-DC boost converter. The proposed control scheme eliminated steady state oscillations around the working point area of the PV panel when there were rapid changes to the irradiance. **Applications/Improvements:** Furthermore, the PV system became more efficient, as proven by the sudden change in the radiation conditions for 0.3 seconds, where approximately 30% of energy could be saved.

Keywords: Controlled Mobility, Energy Efficiency, MANETs

1. Introduction

MANETs are gaining popularity because of its ease deployment. MANET is a decentralized infrastructure less network and the nodes move arbitrarily¹. The nodes communicate either directly or through multi hops. MANETs face various problems like unpredictable topology, increased interference and congestion and limitation of resources like bandwidth and energy due to shared wireless medium and its dynamic nature. The nodules (nodes) contained by an ad hoc system usually rely on batteries (or exhaustive energy sources) for control (power). Since these control sources have a restricted duration, power ease of use is one of the most vital limits for the process of the ad hoc system².

There are different sources of power consumption in a mobile node. Energy is consumed while sending Packet, while receiving a packet, in idle mode and in sleep mode

which occurs when the wireless interface of the Mobile node is turned off. Power utilization can be optimized by employing routing algorithms that avoid nodes with less battery power remaining while trying to minimize the total power consumed in transmitting a packet. In mobile ad hoc network there are three aspects to reduce the power consumption - Power saving at mobile device level, Power saving by controlling transmission level of packet and Power saving by using optimized power routing protocol.

In this study, a heterogeneous framework comprises of conventional nodules and transfer nodules. Mobile relay nodules are used to diminish control utilization in MANETs. Relay nodes are costly, more powerful and their main task is communication with other nodes. Relay nodes have same transmission radii. Nodules consume extra control when it is aggressively sending data packet when compared to packet response and inactive stage.

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Also transmitting the data over extended distances it uses extra power in contrast to small distance transmission.

The heterogeneous MANET cooperates on the power utilized at the relay nodules in turn to diminish the energy utilized at the conventional nodules. The relay nodules can be placed nearer to the conventional nodules to diminish the energy utilized throughout packet broadcasting. In this paper two, methods of the relay deployment problem together with their solutions is presented.

2. Energy Model

In wireless networks, communication takes place over multi hop path. The data flow between starting and ending nodes can be characterized as a set of one step runs among the adjoining nodules³. Let T be the set of traditional nodes and R represents the set of mobile relay nodes. F be the set of all one hop flows. Each flow is represented as $f_i = (s_i, d_i, \lambda_i)$, $i = 1, 2, \dots, N$, where s_i is the source node, d_i is the destination node and λ_i is the data rate. The energy utilized at the source node by a course f_i for each component (unit) time is specified by:

$$E_i = \lambda_i * P_T(\delta_{s_i d_i})$$

Where P_T is the energy required for sending a single fragment (bit) of data among the nodules. Power function P_T is generally specified as $P_T = a + b\delta_{sidi}$, where a and b depends on the uniqueness of the communiqué channel. δ_{sidi} is the distance between the source node and the destination node. Let F_r denote the service set of the spread nodule $r_j \in R$. F_r consists of every part of the courses that are being relayed throughout the nodule r_j and $F_r \subseteq F$. The energy utilized for transmitting the data for each unit of instant at nodule s_i for a run belongs to the check (service) set of the relay nodule s_i specified by:

$$E_i^j = \lambda_i * P_T(\delta_{s_i r_j})$$

2.1 Min-Total Formulation

This formulation aims at minimizing the whole power consumed transversely the conventional nodules in a check set. Here a relay nodule is placed in order to diminish the whole power consumed transversely all the conventional nodules in its check set. The entire power inspired by the conventional nodules in the absence of

relay nodules for data transmission per unit time is given by:

$$E_{tot} = \sum_{f_i \in F} E_i$$

The entire power inspired for data transmission by the conventional nodules in the occurrence of relay nodules is specified by:

$$E_{totr} = \sum_{j \in R} \sum_{f_i \in F_r} E_i^j + \sum_{f_k \in F - F_r} E_k$$

The total transmission energy consumed in a relay enabled network is dependent on the position of the relay nodes, position of the conventional nodules and set of runs (flows) in the system, mutual check set of the relay nodules. The position of relay nodule and the mutual check set of relay nodule are adjusted so that the overall energy consumed across the conventional nodule is minimum. The drawback of this method is that the power utilization across the entity nodules might be irregular.

2.2 Min-Max Formulation

The objective of the Min-Max formulation is to diminish the utmost transmission energy utilized transversely the conventional nodules in the system. The utmost energy utilized by a conventional nodule in the check set F_{rj} is given by:

$$E_{max}^j = \max_{i \in F_{rj}} E_i^j$$

This formulation aims at minimizing E_{jmax} for the service set F_{rj} by finding the most advantageous place of the relay nodule. To accomplish this, the best probable check set for each relay nodule has to be defined and computed.

2.3 Assigning Node Weights

In the above two formulations the relay node is placed nearer to the starting place of the nodules with advanced data flow tariffs. In some cases, the assigning nodule weights recommended with the intention of the relay nodule favours the source node with low residual energy. Some nodes in a network will be utilized often and such nodes should be given more weightage to prolong the lifetime of the network.

These factors are included in the crisis formulation by allocating a load w_i to every course $f_i \in F_{r_j}$. The load w_i is the precedence allocated to the course f_i and is given by:

$$w_i = z_1 * \epsilon_i + z_2 * p_i$$

Where ϵ_i and p_i denotes the lasting power and related precedence of the source nodule s_i equivalent to the course f_i . z_1 and z_2 are the comparative loads allocated to the lasting power and related precedence by the system operator. Min-Total scheme after assigning node weights is given by:

$$E_{totr} = \sum_{j \in R} \sum_{f_i \in F_j} w_i * E_i^j + \sum_{f_k \in F - F_r} w_j E_k$$

The optimization function of Min-Max scheme after assigning node weights is given by:

$$E_{max}^j = \max_{i \in F_{r_j}} w_i * E_i^j$$

3. Static Networks

The solutions for a mobile network are initially we have to express the results for a stationary (static) system and after that proceed. The locating and association of the relay nodule depends on the location of the conventional nodules and the set of energetic courses involving with the system. There are two steps to be carried out to solve the two problem instances: 1. Find the every relay node is set as optimal service. 2. The objective function is based on compute the optimal position of the relay node. The two steps are closely related to each other and point to an iterative clarification as an alternative of a successive advance.

3.1 Placing the Relay Node

Assuming a system with eight standing conventional nodules $\{n_1, n_2, \dots, n_8\}$ and a particular relay nodule r_1 as shown in Figure 1a. Here the pair of energetic back-to-back multihop runs in the system with rates of data λ_1 and λ_2 . The pair of one step runs F in the system can be written as:

$$F = \{(n_1, n_2, \lambda_1), (n_2, n_3, \lambda_1), (n_3, n_4, \lambda_1), (n_4, n_5, \lambda_1), (n_1, n_6, \lambda_2), (n_6, n_7, \lambda_2), (n_7, n_8, \lambda_2)\}$$

For a relay node to serve a flow the corresponding sender-receiver couple be supposed to be inside the communication variety of the relay nodule. From the Figure 1(b). Assuming that the courses (n_3, n_4, λ_1) and (n_7, n_8, λ_2) are served by relay node, after that the desirable check sets of r_1 are $\{(n_3, n_4, \lambda_1), (n_7, n_8, \lambda_2)\}$, $\{(n_3, n_4, \lambda_1)\}$ and $\{(n_7, n_8, \lambda_2)\}$.

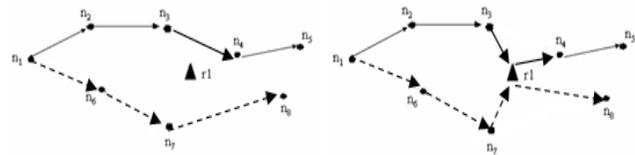


Figure 1. a) Information runs with an exclusive of relay nodule. b) Information runs after arranging a relay nodule.

4. Cellular Networks

In a cellular system, the nodules are extremely active and results in frequent topology changes. The location of traditional nodule continuously changes and the location of the relay nodule be supposed to be calculated for every moment. The mobility epochs is called as the function of the system is modeled more than the predetermined extent instance. The movement and the positioning of the relay nodules depend on the calculated alterations in the system condition above the duration of a time⁴.

4.1 Clarification for Min-Total

The first step is to determine the most favorable check set of the relay and calculate the most advantageous relay's location⁵. Also the communicating node pair of a flow in a check set be supposed to be in the variety (range) of the relay nodule of the epoch duration is throughout.

The relay nodule's check set and movement of the relay node is addressed by the greedy heuristic between cyclic dependency algorithms in Figure 2 This algorithmic step is for each relay nodule in the system at the establishment of each period executed in a distributed fashion⁶. As the first step, all the check desires whose equivalent node pair is not currently in its transmission choice the discards relay node. After that it includes the course by means of the highest load in the demand set to its check set. The course of the relay nodule is known as by means of depends on the calculated location of the nodules in its check set. The relay nodule afterward iteratively efforts to include extra courses to its check set depends on the downward arrangement.

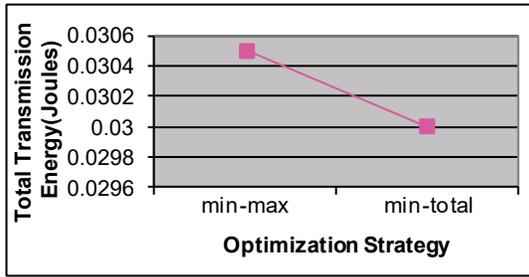


Figure 2. Min-Max vs. Min-Total.

4.2 Clarification for Min-Max

The Min-Total problem is unlike, we adopt a different approach to mobile setup in calculate the Min-Max clarification⁷. During the first step, the check set of the relay nodule includes of each and every one of the runs of the conversing nodule couples are presently in the choice of the relay nodule transmission on behalf of each resource nodule in its check set, the relay nodule subsequently calculates their location at a predetermined integer of occasion time above the itinerary of a time. Nevertheless, as an alternative of calculating a part Min-Max clarification on behalf of all moments, initially think about the nodule positions at each and every one these moments as a dissimilar position and calculate a particular Min-Max clarification for the total period⁸. Depends on this calculated location for the relay nodule, ensure that all the conversing nodule couples are calculated to be in the broadcasting variety of the relay nodule. Suppose if the relay nodule is not in the range, then the course by means of the lowly load is detached as of the check set and then this method repeats in anticipation of every one of the nodules in the check set are calculated to be in the broadcasting variety of the relay nodule or the check set includes of a lone course⁹. Figure 1 shows above described as the Min-Max algorithmic steps for a cellular phone situation. Stipulation the check set includes of a solitary course, the relay is located as near to the resource nodule to save greatest power in the resource nodule.

5. Results and Discussion

The recital of two algorithms (Min-Max and Min-Total) is evaluated.

Figure 2, it is found that; Min-Total algorithm results in improved results of largely power consumption. Min-Max algorithm results in optimization of largest power

inspired by a nodule, but the overall energy consumption is more. Also, when the number of nodes supported by the relay node increases, the total transmission energy decreases as given in Figure 3.

When the density of node is large, equally the two algorithms are executed in a same manner. The differentiation in the algorithm performance is obviously marked while the nodule compactness is fewer¹⁰. If the numeral of relay nodules is improved, the number of courses provided by the relay nodule improves. Accordingly, the communication power absorbed by the conventional nodule reduces ensuing in recital growths. Also with longer epoch duration, the relay node supports limited flows resulting in decreasing performance gains.

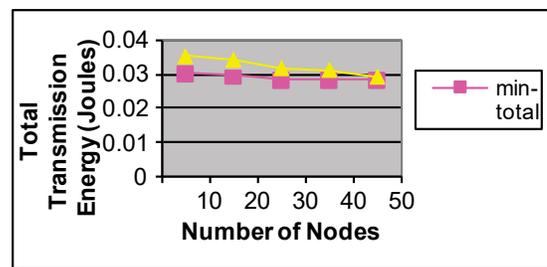


Figure 3. Total transmission energy by varying nodes.

6. Conclusions

In this proposed method, the difficulty of creating relay nodules to preserve control in Mobile Ad hoc Network is studied. Two illustration of the relay exploitation difficulty, simultaneously with their clarifications are presented. Illustration 1, expressed as Min-Total, during the data transmission it diminishes the entire power utilized diagonally all the conventional nodules, whereas in illustration 2, expressed as Min-Max, is used to diminishes the greatest power utilized by a conventional node for the period of data transmission. The proposed framework minimizes the power (energy) utilization at the conventional nodules while compromising on the power possessions at the relay nodules. Simulation results indicate that, arranging a tiny portion of relay nodules by means of the planned structure is not simply produce the outcome within considerable power investments however it can also enhances the entire recital of the system. Combining both Min-Total and Min-Max results in better energy savings.

7. References

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