# Application of Modified Cramer's Rule in Quadrant Interlocking Factorization 

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#### Abstract

Objectives: To show that modified Cramer's rule is better than classical Cramer's rule for solving linear systems in quadrant interlocking factorization or WZ factorization. Methods: The relative residual measurement of modified Cramer's rule was compared with classical Cramer's rule. Furthermore, we apply the rules in WZ factorization and evaluate their matrix norm on AMD and Intel processor. Findings: This study shows that the residual measurements of modified Cramer's rule are $20 \%$ better than Cramer's rule. It also shows that the matrix norm of Cramer's rule in WZ factorization is higher than using modified Cramer's rule in the factorization. Application/improvements: Modified Cramer's rule can be used to solve simple linear system. Applying the modified Cramer's rule in $W Z$ factorization using parallel computer or shared memory multiprocessor networks such as Intel Xeon Phi, Sunway Taihulight or OLCF-4 should be strongly considered.


Keywords: Cramer's Rule, Matrix Norm, Quadrant Interlocking Factorization, WZ Factorization, Z- matrix

## 1. Introduction

Quadrant interlocking factorization (QIF) or $W Z$ factorization produces a $W$-matrix together with a $Z$-matrix (or hourglass matrix) from nonsingular matrix $A^{1,2}$ such that

$$
\begin{equation*}
A=W Z \tag{1}
\end{equation*}
$$

The necessary condition for nonsingular matrix $A=\left[a_{i, t}\right]_{, j,=1}^{n}$ to be a WZ factorization is that the central submatrices $A_{n=m+1}^{c}=\left[a_{i, j}\right]_{i, j+1}^{u m+1}$ are centro-nonsingular ${ }^{3}$. Where n is even order of matrix and $c$ the centered submatrix of $A$, for $\mathrm{m}=1,2, \ldots,\left[\frac{n}{2}\right]$. A matrix which is either a $Z$-matrix or a $W$-matrix is called butterfly matrix. These names are suggested by the shapes of the set of all possible positions for nonzero entries, which are as follows


To obtain the entries of Z-matrix by updating matrix $A$, the column of $W$-matrix must be computed from $2 \times 2$ systems of linear equations in equation 2 using Cramer's rule

$$
\left\{\begin{array}{c}
z_{(m, m)}^{(m-1)} w_{(i, m)}^{m}+z_{(n-m+1, m)}^{(m-1)} w_{(i, n-m+1)}^{m}=-z_{(i, m)}^{(m-1)}  \tag{2}\\
z_{(m, n-m+1)}^{(m-1)} w_{(i, m)}^{m}+z_{(n-m+1, n-m+1)}^{(m-1)} w_{(i, n-m+1)}^{m}=-z_{(i, n-m+1)}^{(m-1)}
\end{array}\right.
$$

to obtain equation (3) as

$$
\begin{equation*}
z_{i, j}^{m}=z_{i, j}^{m-1}+w_{i, m}^{m} z_{m, j}^{m-1}+w_{i, n-m+1}^{m} z_{n-m+1, j}^{m-1} \tag{3}
\end{equation*}
$$

For $i, j=m+1, \ldots, n-m$
For vivid understanding of $W Z$ factorization, readers may see ${ }^{4-11}$ and the references therein. One of the advantages of Cramer's rule is to determine if a system of linear equations is inconsistent or indeterminate because Cramer's rule gives a clear representation of an individual component unconnected to all other components ${ }^{12,13}$. This specific advantage of Cramer's rule is useful in $W Z$ factorization to solve for the $2 \times 2$ linear systems which is the major factor to know if the matrix is centro-nonsingular. The matrix norm of $W Z$ factorization is the Frobenius norm of the matrix given as ${ }^{14}$

$$
\|A\|_{F}=\|A-W Z\|
$$

The Frobenius norm or the Hilbert-Schmidt norm is equivalent to the Euclidean norm on $K^{n \times n} \underline{15}$.

In Section 2, we apply the two modified methods of Cramer's rule proposed by ${ }^{16}$ in $W Z$ factorization. We further evaluate their matrix norm using MATLAB on two distinct physical processors (AMD and Intel processor). The results obtained from the matrix norm can be improved if parallel computer or mesh multiprocessors are used.

## 2. Applying Modified Cramer's Rule in QIF or WZ factorization

### 2.1 Modified Cramer's Rule

Irrespective of its high computational time, Cramer's rule is hypothetically significance for solving systems of linear equations because accurate method to evaluate determinants can make Cramer's rule numerically stable ${ }^{17-19}$. Thus, much deductions have been made on Cramer's rule to solve simple and large-scale systems of linear equations, see for example ${ }^{20-22}$.
Theorem 2.1.1 [Cramer's rule]. Let $A x=c$ be an $n \times n$ system of linear equation and $A$ an $n \times n$ matrix such that $|A| \neq 0$, then the unique solution $x=\left(x_{1}, x_{2}, \ldots, x_{n}\right)^{T}$ to the system of linear equations is given by

$$
\begin{equation*}
x_{i}=\frac{\left|A_{i c}\right|}{|A|} \tag{4}
\end{equation*}
$$

where, $A_{(i \mid c)}$ is the matrix obtained from $A$ by substituting the column vector $c$ to the $i$ th column of $A$, for $i=1,2, \ldots, n$
$\mathrm{In}^{16}$, the following corollaries were obtained from Theorem 2.1.1 which are based on one of the properties of determinant.

Corollary 2.1.1 Let $A x=c$ be an $n \times n$ system of linear equation and $A$ an $n \times n$ matrix such that $|A| \neq 0$, then the unique solution $x=\left(x_{1}, x_{2}, \ldots, x_{n}\right)^{\mathrm{T}}$ to the system of linear equations is given by

$$
\begin{equation*}
x_{i}=\frac{\left|A_{i+c}\right|}{|A|}-1 \tag{5}
\end{equation*}
$$

where, $A_{(i+c)}$ is the matrix obtained from $A$ by adding the column vector $c$ to the $i$ th column of $A$, for $i=1,2, \ldots, n$

Corollary 2.1.2. Let $A x=c$ be an $n \times n$ system of linear equation and $A$ an $n \times n$ matrix such that $|A| \neq 0$, then the unique solution $x=\left(x_{1}, x_{2}, \ldots, x_{n}\right)^{\mathrm{T}}$ to the system of linear equations is given by

$$
\begin{equation*}
x_{i}=1-\frac{\left|A_{1-c}\right|}{|A|} \tag{6}
\end{equation*}
$$

where, $A_{(i-c)}$ is the matrix obtained from $A$ by subtracting the column vector $c$ from the $i$ th column of $A$, for $i=1,2, \ldots, n$.

The authors proved that Corollary 2.1.1 and Corollary 2.1.2 are equal to classical Cramer's rule. Now, we further the findings by providing the algorithm for the above corollaries in Figure 1.

Even though the modified methods of Cramer's rule use extra $n$ arithmetic operations for every solution of the linear systems, their computational time is like Cramer's rule. Therefore, the modified Cramer's rule may nowhere better than other direct methods (such as GE or LU decomposition) in terms of computational time complexity or numerical stability for higher linear systems. Though, our main objective is to apply the modified Cramer's rule in QIF but the methods can be compared with Cramer's rule in terms of their relative residual error, $\frac{\|A \bar{x}-c\|}{\|A\|\|\|}$, which is given in Table 1.

Figure 2 shows that the relative residual of Method I and Method II are almost similar and are better than residual of Cramer's rule. Among the three algorithms, Method II shows to be better than Cramer's rule and Method I. Our algorithms, for Method I and Method II, are about $20 \%$ better than Cramer's rule.

### 2.2 Application of Modified Cramer's Rule

Due to the lack of parallel computer, the MATLAB codes application of this paper are limited to AMD and Intel processor with standard hardware in Table 2.

```
Algorithm 1 Modified Cramer's rule algorithm
    procedure Modified Cramer's rule
            \(A \leftarrow n \times n\) coefficient matrix
            \(c \leftarrow\) column vector
            \(x_{i} \leftarrow\) solutions of linear system
            for i doln
                \(K \leftarrow\) the \(i\) th column of \(A \pm c\)
                \(\operatorname{det}(A) \neq 0 \leftarrow\) determinant of \(A\)
                \(\operatorname{det}(K) \leftarrow\) determinant of \(K\)
                if \(K=S+c\) then
                    \(x_{i} \leftarrow 1-(\operatorname{det}(K) / \operatorname{det}(A))\)
            else
                \(x_{i} \leftarrow(\operatorname{det}(K) / \operatorname{det}(A))-1\)
                end if
            end for
            return \(x_{i}\)
    end proced ure
```

Figure 1. Algorithm of modified Cramer's rule.


Figure 2. Residual measurement of Cramer's rule, Method I and Method II.

To apply the above corollaries in $W Z$ factorization algorithm, we reconstruct equation (2) as

$$
\underbrace{\left[\begin{array}{cc}
z_{m, m}^{m-1} & z_{n-m+1, m}^{m-1}  \tag{7}\\
z_{m, n-m+1}^{m-1} & z_{n-m+1, n-m+1}^{m-1}
\end{array}\right]}_{A} \underbrace{\left[\begin{array}{c}
w_{i, m}^{m} \\
w_{i, n-m+1}^{m}
\end{array}\right]}_{x}=\underbrace{-z_{i, m}^{m-1}}_{c} \begin{gathered}
-z_{i, n-m+1}^{m-1}
\end{gathered}
$$

Now, we apply Corollary 2.1.1 to compute $w_{i, m}^{m}$ and $w_{i, n-m+1}^{m}$ from equation (7) to obtain

$$
\begin{equation*}
w_{i, m}^{m}=\frac{\left|\left[w_{i, m}^{m}\right]_{i+c}\right|}{|z|}-1 \text { and } w_{i, n-m+1}^{m}=\frac{\left|\left[w_{i, n-m+1}^{m}\right]_{i+c}\right|}{|z|}-1 \tag{8}
\end{equation*}
$$

where,

$$
\begin{aligned}
& |z|=-z_{n-m+1, m}^{m-1} z_{m, n-m+1}^{m-1}+z_{n-m+1, n-m+1}^{m-1} z_{m, m}^{m-1} \neq 0 \\
& \left|\left[w_{i, m}^{m}\right]_{i+c}\right|=z_{n-m+1, m}^{m-1} z_{i, n-m+1}^{m-1}-z_{n-m+1, n-m+1}^{m-1} z_{i, m}^{m-1}
\end{aligned}
$$

Table 1. Relative residual measurements of Cramer's rule, Method I and Method II

| Matrix size <br> $(\mathrm{N})$ | Cramer | Method | Method |
| :--- | :--- | :--- | :--- |
| $2 \times 2$ | $9.93 \mathrm{E}-17$ | $6.53 \mathrm{E}-17$ | $6.71 \mathrm{E}-17$ |
| $3 \times 3$ | $9.97 \mathrm{E}-17$ | $6.65 \mathrm{E}-17$ | $5.93 \mathrm{E}-17$ |
| $4 \times 4$ | $4.41 \mathrm{E}-16$ | $3.32 \mathrm{E}-16$ | $3.28 \mathrm{E}-16$ |
| $5 \times 5$ | $2.19 \mathrm{E}-16$ | $1.44 \mathrm{E}-16$ | $1.35 \mathrm{E}-16$ |
| $6 \times 6$ | $2.44 \mathrm{E}-16$ | $1.31 \mathrm{E}-16$ | $1.23 \mathrm{E}-16$ |
| $7 \times 7$ | $5.61 \mathrm{E}-16$ | $4.27 \mathrm{E}-16$ | $4.18 \mathrm{E}-16$ |
| $8 \times 8$ | $4.19 \mathrm{E}-16$ | $3.42 \mathrm{E}-16$ | $3.39 \mathrm{E}-16$ |
| $9 \times 9$ | $7.54 \mathrm{E}-16$ | $6.59 \mathrm{E}-16$ | $6.50 \mathrm{E}-16$ |
| $10 \times 10$ | $7.01 \mathrm{E}-16$ | $6.65 \mathrm{E}-16$ | $6.58 \mathrm{E}-16$ |
| $11 \times 11$ | $9.32 \mathrm{E}-16$ | $8.41 \mathrm{E}-16$ | $8.38 \mathrm{E}-16$ |
| $12 \times 12$ | $8.81 \mathrm{E}-16$ | $8.09 \mathrm{E}-16$ | $8.13 \mathrm{E}-16$ |
| $13 \times 13$ | $9.75 \mathrm{E}-16$ | $8.84 \mathrm{E}-16$ | $8.75 \mathrm{E}-16$ |
| $14 \times 14$ | $9.01 \mathrm{E}-16$ | $8.43 \mathrm{E}-16$ | $8.37 \mathrm{E}-16$ |

Table 2. Hardware specifications

| CPU | Memory |
| :--- | :--- |
| AMD APU Core A10-9600P <br> 2.4 GHz | 8 GB |
| Intel Core i7-4600U 2.1 GHz | 8 GB |

$$
-z_{n-m+1, m}^{m-1} z_{m, n-m+1}^{m-1}+z_{n-m+1, n-m+1}^{m-1} z_{m, m}^{m-1}
$$

$$
\begin{gathered}
\left|\left[w_{i, n-m+1}^{m}\right]_{i+c}\right|=z_{m, n-m+1}^{m-1} z_{i, m}^{m-1}-z_{i, n-m+1}^{m-1} z_{m, m}^{m-1} \\
\quad-z_{n-m+1, m}^{m-1} z_{m, n-m+1}^{m-1}+z_{n-m+1, n-m+1}^{m-1} z_{m, m}^{m-1}
\end{gathered}
$$

The factorization obtained from using Corollary 2.1.1 is called $W^{1} Z^{1}$ factorization and its MATLAB code is given in Figure 3.

More so, if we apply Corollary 2.1.2 to compute $\mathrm{w}_{\mathrm{i}, \mathrm{m}}^{\mathrm{m}}$ and $w_{i, n-m+1}^{m}$ from equation (7), we will have

$$
\begin{equation*}
w_{i, m}^{m}=1-\frac{\left|\left[w_{i, m}^{m}\right]_{i-c}\right|}{|z|} \text { and } w_{i, n-m+1}^{m}=1-\frac{\left|\left[w_{i, n-m+1}^{m}\right]_{i+c}\right|}{|z|} \tag{9}
\end{equation*}
$$

where,

$$
|z|=-z_{n-m+1, m}^{m-1} z_{m, n-m+1}^{m-1}+z_{n-m+1, n-m+1}^{m-1} z_{m, m}^{m-1} \neq 0
$$

```
function \(Z=W 1 Z 1\) factorization
    \% step of elimination - from \(A\) to \(Z\)
    \(\mathrm{A}=\) input ('matrix \(\mathrm{A}=\) ');
    \(\mathrm{n}=\operatorname{size}(\mathrm{A}, 1)\);
    \(\mathrm{W}=\operatorname{zeros}(\mathrm{n})\);
    for \(\mathrm{k}=1\) : \(\operatorname{ceil}((\mathrm{n}-1) / 2)\)
        \(\mathrm{k} 2=\mathrm{n}-\mathrm{k}+1\);
        \(\operatorname{det}=\mathrm{A}(\mathrm{k}, \mathrm{k}) * \mathrm{~A}(\mathrm{k} 2, \mathrm{k} 2)-\mathrm{A}(\mathrm{k} 2, \mathrm{k}) * \mathrm{~A}(\mathrm{k}, \mathrm{k} 2)\);
        if \(\operatorname{det}=0\)
            exitflag \(=0\);
            for \(\mathrm{il}=\mathrm{k}: \mathrm{k} 2\)
                for \(\mathrm{i} 2=\mathrm{i} 1: \mathrm{k} 2\)
                \(\operatorname{det}=A(i 1, k) * A(i 2, k 2)-A(i 2, k) * A(i 1, k 2)\);
                if \(\operatorname{det}^{2}=0\)
                    \(\mathrm{tmp}=\mathrm{A}(\mathrm{i} 1, \mathrm{k}: \mathrm{k} 2) ;\)
                        \(\mathrm{A}(\mathrm{i} 1, \mathrm{k}: \mathrm{k} 2)=\mathrm{A}(\mathrm{k}, \mathrm{k}: \mathrm{k} 2)\);
                        \(\mathrm{A}(\mathrm{k}, \mathrm{k}: \mathrm{k} 2)=\mathrm{tmp} ;\)
                    \(\mathrm{tmp}=\mathrm{A}(\mathrm{i} 2, \mathrm{k}: \mathrm{k} 2)\);
                        \(\mathrm{A}(\mathrm{i} 2, \mathrm{k}: \mathrm{k} 2)=\mathrm{A}(\mathrm{k} 2, \mathrm{k}: \mathrm{k} 2) ;\)
                        \(\mathrm{A}(\mathrm{k} 2, \mathrm{k}: \mathrm{k} 2)=\mathrm{tmp} ;\)
                        exitflag = 1;
                        break
                end
            end
        end
        if exitflag \(=0\)
            \(\mathrm{Z}=\mathrm{A}\);
            return
        end
    end
    \% finding elements of W
    \(W(\mathrm{k}+1: \mathrm{k} 2-1, \mathrm{k})=((\mathrm{A}(\mathrm{k} 2, \mathrm{k}) * \mathrm{~A}(\mathrm{k}+1: \mathrm{k} 2-1, \mathrm{k} 2)-\mathrm{A}(\mathrm{k} 2, \mathrm{k} 2) * \mathrm{~A}(\mathrm{k}+1: \mathrm{k} 2-1, \mathrm{k})-\mathrm{A}(\mathrm{k} 2, \mathrm{k}) * \mathrm{~A}(\)
        \(\mathrm{k}, \mathrm{k} 2)+\mathrm{A}(\mathrm{k} 2, \mathrm{k} 2) * \mathrm{~A}(\mathrm{k}, \mathrm{k})) /\) det \()-1\);
    \(W(k+1: k 2-1, k 2)=((A(k, k 2) * A(k+1: k 2-1, k)-A(k+1: k 2-1, k 2) * A(k, k)-A(k 2, k) * A(k\)
        \(, \mathrm{k} 2)+\mathrm{A}(\mathrm{k} 2, \mathrm{k} 2) * \mathrm{~A}(\mathrm{k}, \mathrm{k})) / \operatorname{det})-1\);
    \% updating A
    \(\mathrm{A}(\mathrm{k}+1: \mathrm{k} 2-1, \mathrm{k}+1: \mathrm{k} 2-1)=\mathrm{A}(\mathrm{k}+1: \mathrm{k} 2-1, \mathrm{k}+1: \mathrm{k} 2-1)+\mathrm{W}(\mathrm{k}+1: \mathrm{k} 2-1, \mathrm{k}) * \mathrm{~A}(\mathrm{k}, \mathrm{k}+1: \mathrm{k} 2\)
        \(-1)+W(k+1: k 2-1, k 2) * A(k 2, k+1: k 2-1)\);
    \(\mathrm{Z}=\mathrm{A}\);
end
```

Figure 3. MATLAB code of $W^{1} Z^{I}$ factorization.

```
    Finding elements of W 
    W(k+1:k2-1,k)=1-((A(k2,k2)*A(k+1:k2-1
    W(k+1:k2-1,k2)=1-((A(k,k)*A(k+1:k2-1,k2)-A(k,k2)*A(k+1:k2-1,k)-A(k2,k
            )*A(k,k2)+A(k2,k2)*A(k,k))/det);
```

Figure 4. MATLAB code of $W^{2} Z^{2}$ factorization.

```
%finding elements of W
    W(k+1:k2-1,k)=(A(k2,k)*A(k+1:k2-1,k2)-A(k2,k2)*A(k+1:k2-1,k))/det;
    W(k+1:k2-1,k2)=(A(k,k2)*A(k+1:k2-1,k)-A(k,k)*A(k+1:k2-1,k2))/det;
```

Figure 5. MATLAB code of $W Z$ factorization.


Figure 6. Norm of $W Z, W^{1} Z^{1}$ and $W^{2} Z^{2}$ and factorization on Intel processor.

Table 3. Norms of $W Z, W^{1} Z^{1}$ and $W^{2} Z^{2}$ factorization on Intel processor

| Matrix size (N) | \||A-WZ|| | $\left\|\left\|A-W^{2} Z^{2}\right\|\right\|$ | \||A- $W^{2} Z^{2}\| \|$ |
| :---: | :---: | :---: | :---: |
| $10 \times 10$ | $0.47 \mathrm{E}-7$ | 0.18E-7 | $0.09 \mathrm{E}-7$ |
| $20 \times 20$ | 0.88E-7 | 0.53E-7 | 0.42E-7 |
| $30 \times 30$ | $1.47 \mathrm{E}-7$ | $1.13 \mathrm{E}-7$ | $1.04 \mathrm{E}-7$ |
| $40 \times 40$ | $1.95 \mathrm{E}-7$ | $1.51 \mathrm{E}-7$ | $1.44 \mathrm{E}-7$ |
| $50 \times 50$ | $2.58 \mathrm{E}-7$ | 2.17E-7 | $2.08 \mathrm{E}-7$ |
| $60 \times 60$ | $2.93 \mathrm{E}-7$ | $2.61 \mathrm{E}-7$ | $2.52 \mathrm{E}-7$ |
| $70 \times 70$ | $3.47 \mathrm{E}-7$ | $3.12 \mathrm{E}-7$ | $3.02 \mathrm{E}-7$ |
| $80 \times 80$ | $3.89 \mathrm{E}-7$ | $3.58 \mathrm{E}-7$ | $3.49 \mathrm{E}-7$ |
| $90 \times 90$ | 4.47E-7 | 4.17E-7 | $4.05 \mathrm{E}-7$ |
| $100 \times 100$ | $4.99 \mathrm{E}-7$ | 4.73E-7 | $4.59 \mathrm{E}-7$ |
| $110 \times 110$ | $5.49 \mathrm{E}-7$ | 5.14E-7 | 5.03E-7 |
| $120 \times 120$ | $5.96 \mathrm{E}-7$ | $5.58 \mathrm{E}-7$ | $5.41 \mathrm{E}-7$ |
| $130 \times 130$ | $6.49 \mathrm{E}-7$ | $6.11 \mathrm{E}-7$ | $6.02 \mathrm{E}-7$ |
| $140 \times 140$ | $6.98 \mathrm{E}-7$ | $6.60 \mathrm{E}-7$ | $6.52 \mathrm{E}-7$ |
| $150 \times 150$ | 7.47E-7 | 7.10E-7 | 7.01E-7 |
| $160 \times 160$ | $7.83 \mathrm{E}-7$ | 7.57E-7 | 7.43E-7 |
| $170 \times 170$ | $8.55 \mathrm{E}-7$ | $8.25 \mathrm{E}-7$ | 8.16E-7 |
| $180 \times 180$ | 8.97E-7 | 8.52E-7 | 8.46E-7 |
| $190 \times 190$ | $9.58 \mathrm{E}-7$ | $9.18 \mathrm{E}-7$ | $9.03 \mathrm{E}-7$ |
| $200 \times 200$ | $9.99 \mathrm{E}-7$ | $9.68 \mathrm{E}-7$ | $9.57 \mathrm{E}-7$ |

$$
\begin{gathered}
\left|\left[w_{i, m}^{m}\right]_{i-c}\right|=z_{n-m+1, n-m+1}^{m-1} z_{i, m}^{m-1}+z_{n-m+1, n-m+1}^{m-1} z_{m, m}^{m-1} \\
-z_{n-m+1, m}^{m-1} z_{i, n-m+1}^{m-1}-z_{n-m+1, m}^{m-1} z_{m, n-m+1}^{m-1} \\
\left|\left[w_{i, n-m+1}^{m}\right]_{i-c}\right|=z_{i, n-m+1}^{m-1} z_{m, m}^{m-1}+z_{n-m+1, n-m+1}^{m-1} z_{m, m}^{m-1} \\
\quad-z_{m, n-m+1}^{m-1} z_{i,}^{m-1}-z_{n-m+1, m}^{m-1} z_{m, n-m+1}^{m-1}
\end{gathered}
$$

We refer the factorization obtained from using Corollary 2.1.2 as $W^{2} Z^{2}$ factorization and its MATLAB


Figure 7. Norm of $W Z, W^{1} Z^{1}$ and $W^{2} Z^{2}$ and factorization on AMD processor.
code is given in Figure 4, where we replaced line 32 and line 33 in Figure 3 with line 2 and line 3 of Figure 4.
$w_{i, n-m+1}^{m}$ In all, if we apply Theorem 2.1.1 to compute $\mathrm{W}_{\mathrm{i}, \mathrm{m}}^{\mathrm{m}}$ and

$$
\begin{equation*}
w_{i, m}^{m}=\frac{\left|\left[w_{i, m}^{m}\right]_{i \mid c}\right|}{|z|} \text { and } w_{i, n-m+1}^{m}=\frac{\left|\left[w_{i, n-m+1}^{m}\right]_{i \mid c}\right|}{|z|} \tag{10}
\end{equation*}
$$

where,

$$
\begin{aligned}
& |z|=-z_{n-m+1, m}^{m-1} z_{m, n-m+1}^{m-1}+z_{n-m+1, n-m+1}^{m-1} z_{m, m}^{m-1} \neq 0 \\
& \left|\left[w_{i, m}^{m}\right]_{i \mid c}\right|=z_{n-m+1, m}^{m-1} z_{i, n-m+1}^{m-1}-z_{n-m+1, n-m+1}^{m-1} z_{i, m}^{m-1} \\
& \left|\left[w_{i, n-m+1}^{m}\right]_{i+c}\right|=z_{m, n-m+1}^{m-1} z_{i, m}^{m-1}-z_{i, n-m+1}^{m-1} z_{m, m}^{m-1}
\end{aligned}
$$

We then refer the factorization obtained from using Theorem 2.1.1 as WZ factorization and its MATLAB code is given in Figure 5, where we replaced line 32 and line 33 in Figure 3 with line 2 and line 3 of Figure 5.

Next, we compute the Frobenius norm of $W Z, W^{1} Z^{1}$ and $W^{2} Z^{2}$ factorization on AMD and Intel processor and the results are recorded in Table 3 and Table 4 respectively.

In Figures 6 and in 7, the norm of $W^{1} Z^{1}$ and $W^{2} Z^{2}$ factorization are relatively close, but they are better than the norm of $W Z$ factorization on both Intel and AMD processor. Due to round off error in Cramer's rule, the norm of $W Z$ factorization is higher. We can deduce that on Intel processor $W Z$ factorization execution has better sequential algorithms than AMD processor. Applying the modified Cramer's rule in WZ factorization using parallel computer or mesh multiprocessors with better FLOPS on shared memory such as Intel Xeon Phi or OLCF-4 is passionately advocated.

Table 4. Norms of $W Z, W^{1} Z^{1}$ and $W^{2} Z^{2}$ factorization on AMD processor

| Matrix size (N) | \||A-WZ|| | \||A- $W^{1} Z^{1}\| \|$ | \||A- $W^{2} Z^{2} \\|$ |
| :---: | :---: | :---: | :---: |
| $10 \times 10$ | 0.66E-7 | $0.44 \mathrm{E}-7$ | $0.34 \mathrm{E}-7$ |
| $20 \times 20$ | 0.83E-7 | 0.53E-7 | $0.45 \mathrm{E}-7$ |
| $30 \times 30$ | 0.96E-7 | $0.61 \mathrm{E}-7$ | $0.53 \mathrm{E}-7$ |
| $40 \times 40$ | 1.07E-7 | 0.81E-7 | 0.69E-7 |
| $50 \times 50$ | 1.13E-7 | 0.92E-7 | 0.83E-7 |
| $60 \times 60$ | $1.36 \mathrm{E}-7$ | $1.11 \mathrm{E}-7$ | $1.03 \mathrm{E}-7$ |
| $70 \times 70$ | $1.63 \mathrm{E}-7$ | $1.39 \mathrm{E}-7$ | $1.24 \mathrm{E}-7$ |
| $80 \times 80$ | 2.02E-7 | $1.78 \mathrm{E}-7$ | $1.64 \mathrm{E}-7$ |
| $90 \times 90$ | 2.42E-7 | $2.16 \mathrm{E}-7$ | $2.03 \mathrm{E}-7$ |
| $100 \times 100$ | 2.97E-7 | $2.62 \mathrm{E}-7$ | $2.49 \mathrm{E}-7$ |
| $110 \times 110$ | 3.53E-7 | $3.24 \mathrm{E}-7$ | $3.13 \mathrm{E}-7$ |
| $120 \times 120$ | 4.29E-7 | $4.02 \mathrm{E}-7$ | $3.93 \mathrm{E}-7$ |
| $130 \times 130$ | 5.15E-7 | $4.79 \mathrm{E}-7$ | $4.67 \mathrm{E}-7$ |
| $140 \times 140$ | 6.08E-7 | 5.76E-7 | $5.59 \mathrm{E}-7$ |
| $150 \times 150$ | 7.08E-7 | $6.83 \mathrm{E}-7$ | $6.71 \mathrm{E}-7$ |
| $160 \times 160$ | 8.09E-7 | 7.77E-7 | 7.63E-7 |
| $170 \times 170$ | 9.16E-7 | 8.74E-7 | 8.56E-7 |
| $180 \times 180$ | 10.35E-7 | $9.92 \mathrm{E}-7$ | $9.75 \mathrm{E}-7$ |
| $190 \times 190$ | $11.42 \mathrm{E}-7$ | 10.98E-7 | 10.84E-7 |
| $200 \times 200$ | $12.69 \mathrm{E}-7$ | 12.32E-7 | 12.25E-7 |

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